

Entropy-Constrained Gaussian Channel Capacity

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SiA group meeting

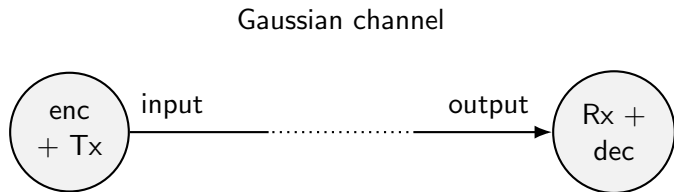
Outline

- 1 Entropy-constrained Gaussian channel
- 2 High-SNR asymptotics
- 3 Low-SNR asymptotics via moment problem

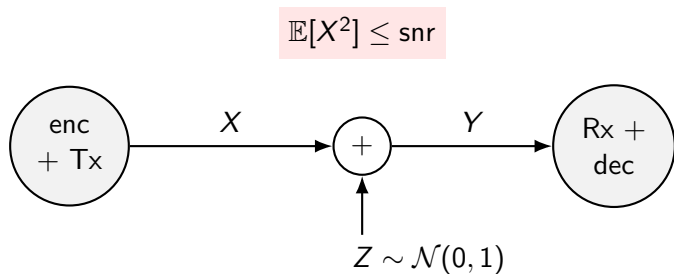
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Gaussian channel



Gaussian channel

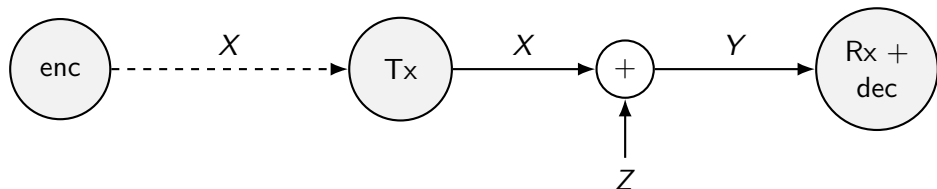


$$C(\text{snr}) = \max_{X: \mathbb{E}[X^2] \leq \text{snr}} I(X; X + Z) = \frac{1}{2} \log(1 + \text{snr})$$

Gaussian channel

finite-capacity noiseless link

$$\mathbb{E}[X^2] \leq \text{snr}$$



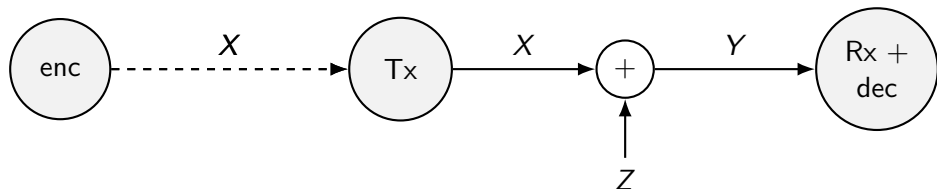
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Entropy-constrained Gaussian channel

finite-capacity noiseless link

$$H(X) \leq h$$

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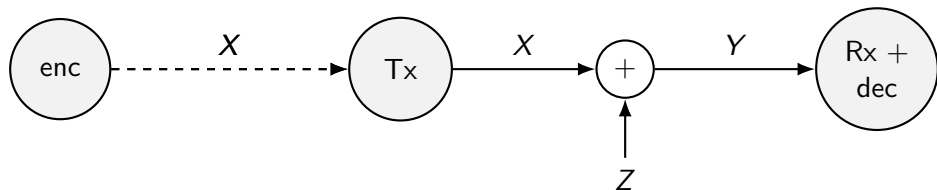
$$C_H(h, \text{snr}) = \max_{\substack{X: \mathbb{E}[X^2] \leq \text{snr} \\ H(X) \leq h}} I(X; X + Z)$$

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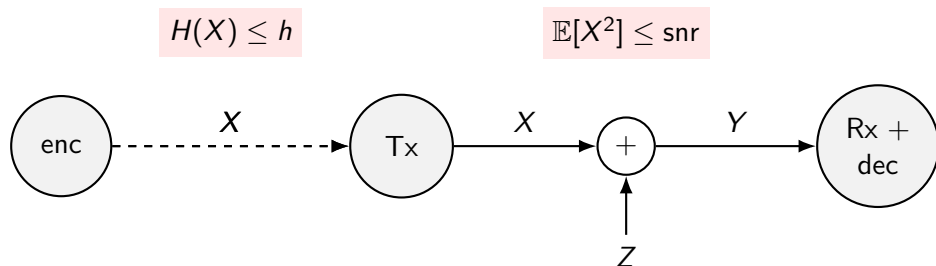


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Entropy-constrained Gaussian channel

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Approximation perspective

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- wlog let $\mathbb{E}[X] = 0$, $\mathbb{E}[X^2] = 1$; $G \sim \mathcal{N}(0, 1)$ independent of Z

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- optimal X at h : discrete X with $H(X) \leq h$ that is closest to $\mathcal{N}(0, 1 + \text{snr})$ after “Gaussian smoothing”

Estimation perspective

Estimation perspective

- MMSE of estimating X from $Y = \sqrt{\text{snr}}X + Z$:

$$\text{mmse}(X, \text{snr}) = \mathbb{E} [(X - \mathbb{E}[X | Y])^2]$$

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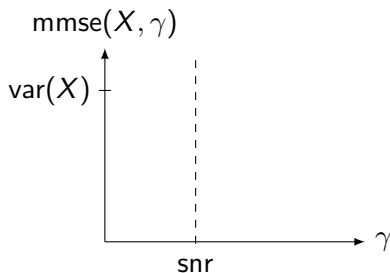
- I-MMSE relationship: $I(X, \text{snr}) = \frac{1}{2} \int_0^{\text{snr}} \text{mmse}(X, \gamma) d\gamma, \quad H(X) = I(X, \infty)$

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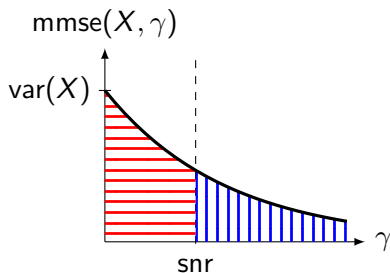


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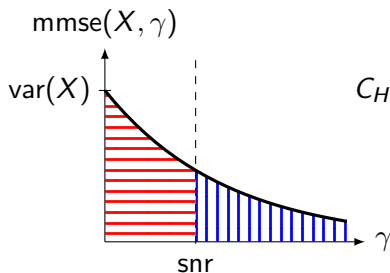


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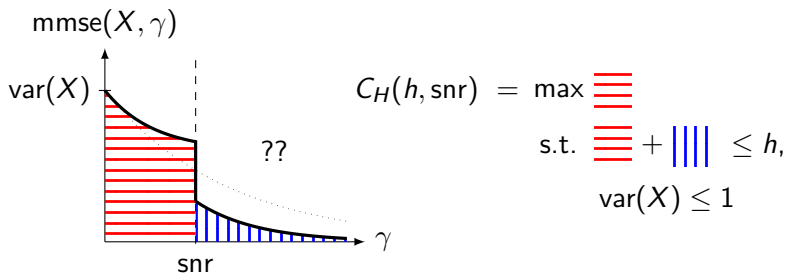
$$C_H(h, \text{snr}) = \max \text{ (red hatched area) } \\ \text{s.t. } \text{ (red hatched area) } + \text{ (blue hatched area) } \leq h, \\ \text{var}(X) \leq 1$$

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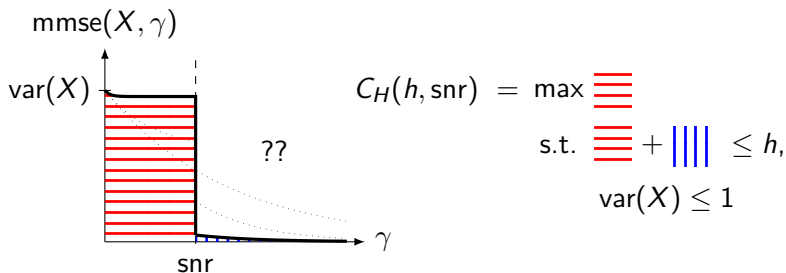


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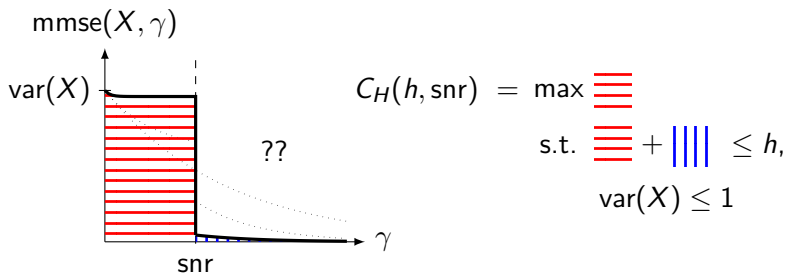


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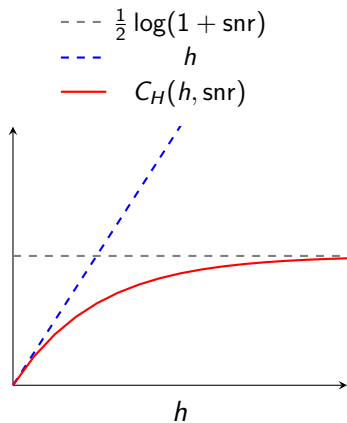
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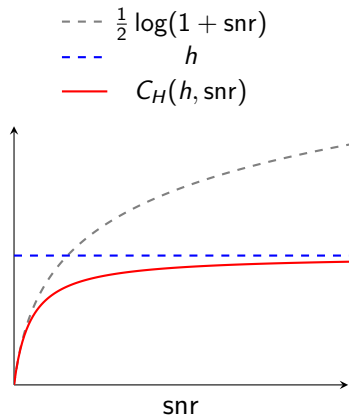
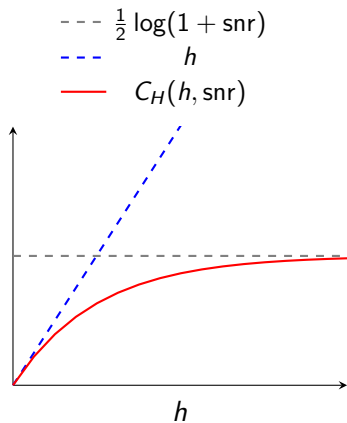
- optimal X at snr : indistinguishable at $\text{SNR} < \text{snr}$, distinguishable at $\text{SNR} > \text{snr}$

C_H as a function of h and snr

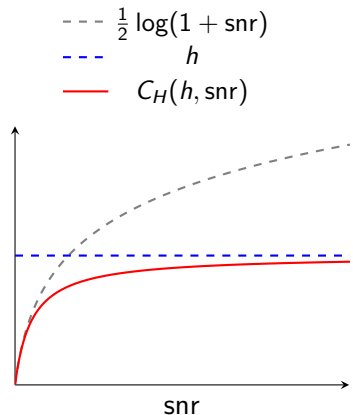
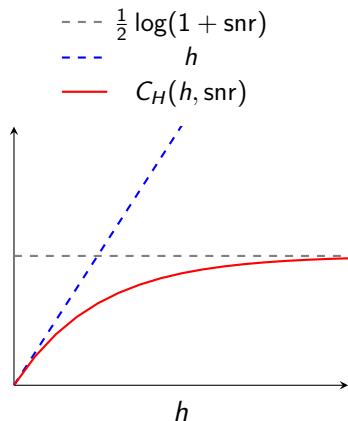
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we look at asymptotic behaviour of C_H as $\text{snr} \rightarrow 0, \infty$

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High-SNR asymptotics

- $\text{snr} \rightarrow \infty : C_H(h, \text{snr}) \rightarrow h$, so gap of interest $= h - C_H(h, \text{snr}) = H(X^*) - I(X^*, \text{snr})$

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Lemma

For any X with $\mathbb{E}[X^2] = 1$ and finite $H(X)$, and $Z \sim \mathcal{N}(0, 1)$ independent of X , we have

$$H(X \mid \sqrt{\text{snr}}X + Z) = \exp \left(-\text{snr} \frac{d_{\min}(X)^2}{8} + o(\text{snr}) \right).$$

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for any other X , there exists snr_0 such that $I(X_h, \text{snr}) > I(X, \text{snr})$ for all $\text{snr} > \text{snr}_0$
- need stronger $X^* \xrightarrow{d} X_h$ as $\text{snr} \rightarrow \infty$ to characterize **gap** in terms of $d_{\min}(X_h) = \beta$

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- $\mathbb{E}[X^{2n}] < \infty$: $I(X, \text{snr}) = \sum_{i=1}^{n-1} a_{i,X} \text{snr}^i + r_{n,X} \text{snr}^n$

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- $k_h :=$ maximum number of moments of G matched by X with $H(X) \leq h$
- $C(\text{snr}) - C_H(h, \text{snr}) = O(\text{snr}^{k_h+1})$

Aside: computing integrals numerically

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- given continuous $W \sim f_W$ and $[C]$ constraints on discrete X , how large can $k_{[C]}$ be?

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A: iff $H_n(s_1, \dots, s_{2n}) = \begin{pmatrix} 1 & s_1 & \dots & s_n \\ s_1 & s_2 & \dots & s_{n+1} \\ \vdots & \vdots & \ddots & \vdots \\ s_n & s_{n+1} & \dots & s_{2n} \end{pmatrix} \succeq 0$ for $n = 1, 2, 3, \dots$

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finite support: at most $\lfloor k/2 \rfloor + 1$ atoms (if it exists)

Gauss–Hermite quadrature

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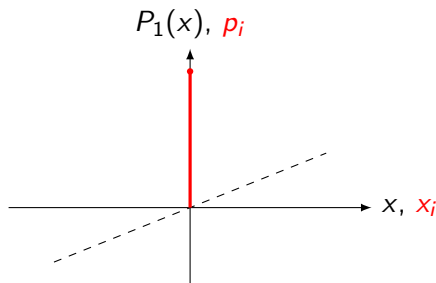
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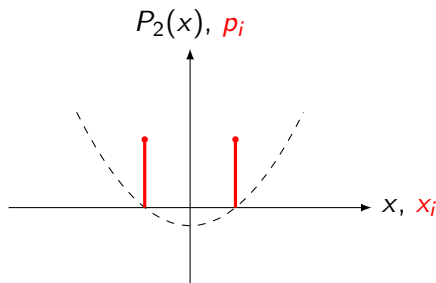


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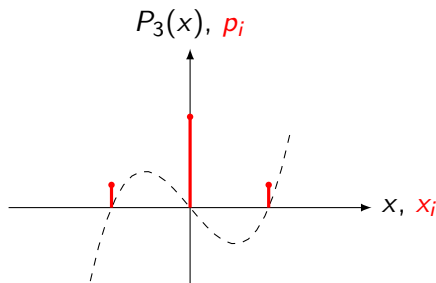
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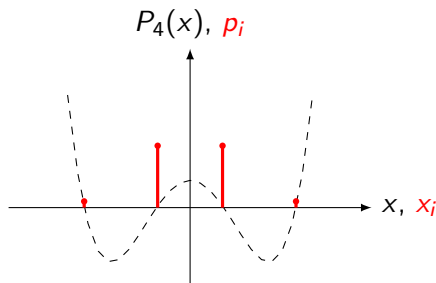


$$\begin{aligned}\mathbb{E}[X] &= 0 & \mathbb{E}[X^2] &= 1 \\ \mathbb{E}[X^3] &= 0 & \mathbb{E}[X^4] &= 3 \\ \mathbb{E}[X^5] &= 0 & \mathbb{E}[X^6] &\neq 15\end{aligned}$$

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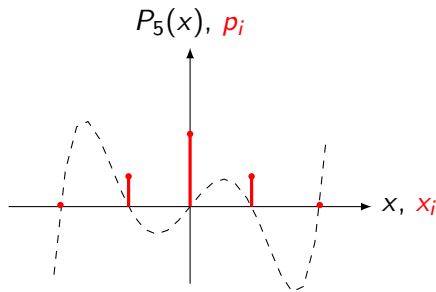
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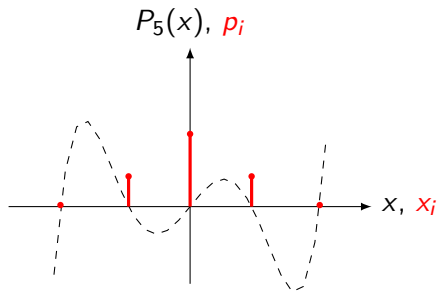
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- optimal trade-off of atoms to moments matched, but $H(X_m^Q) \approx \frac{1}{2} \log m \rightarrow \infty$

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Corollary

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Corollary

$\eta_G = \frac{1}{3}$, so for $h < h_2(\frac{1}{3})$, as $\text{snr} \rightarrow 0$, $C(\text{snr}) - C_H(h, \text{snr}) = O(\text{snr}^4)$.

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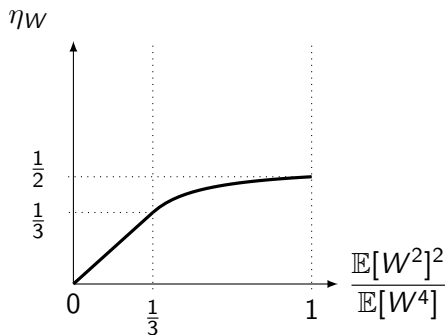
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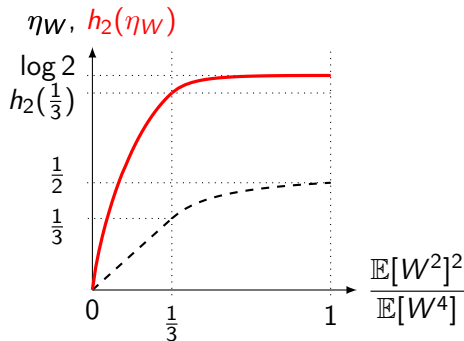
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- proof: Taylor expansion, k^{th} term depends on first k moments, can match only 3 moments

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Thank you!