

Trivial is (Sometimes) Best: Distributed Hypothesis Testing via Linear Codes

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joint work with Robinson Cung and Emre Telatar

EPFL



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LIDS student conference

Binary hypothesis testing

- $\mathcal{H} \in \{0, 1\} : Z \sim P_Z^{\mathcal{H}}$

Binary hypothesis testing

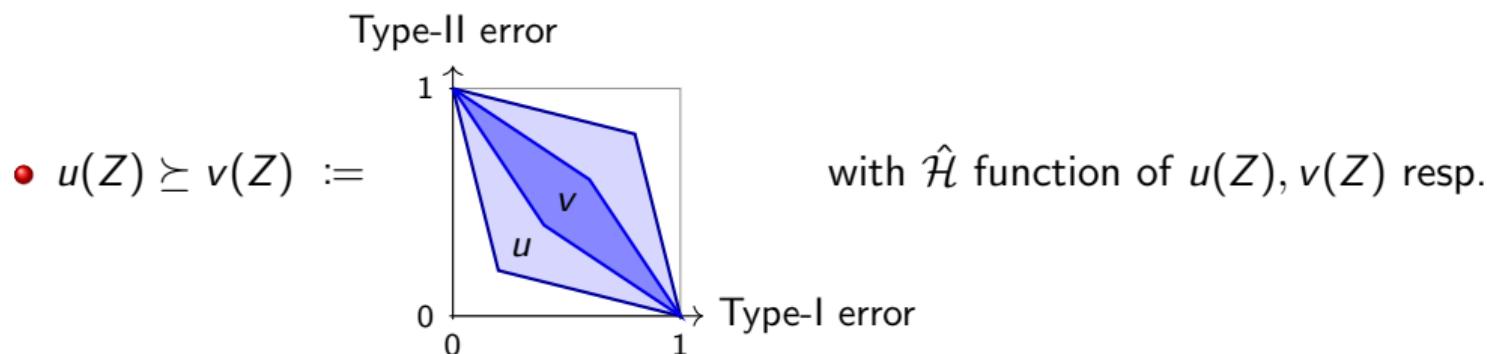
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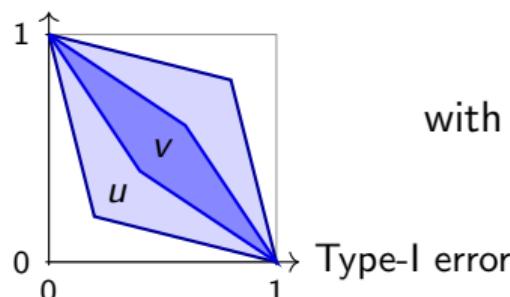


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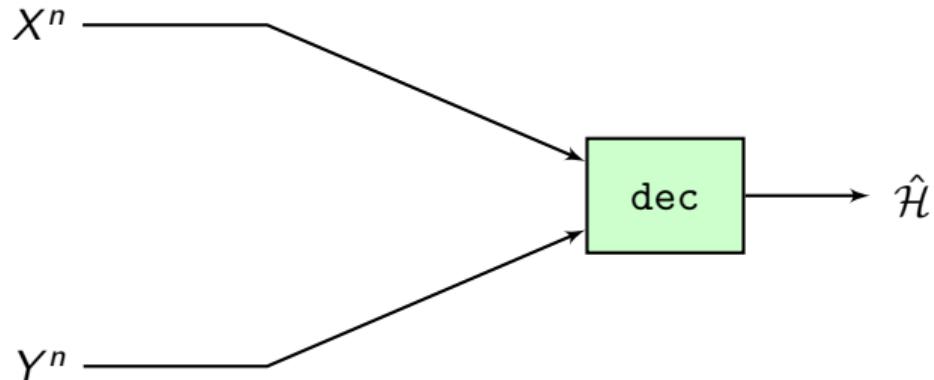
- $u(Z) \succeq v(Z) :=$ with $\hat{\mathcal{H}}$ function of $u(Z), v(Z)$ resp.



- $u(Z) \succeq v(Z) \iff$ there exists $P_{V|U}$ such that $P_{u(Z)}^{\mathcal{H}} \xrightarrow{P_{V|U}} P_{v(Z)}^{\mathcal{H}}$ for both $\mathcal{H} = 0, 1$

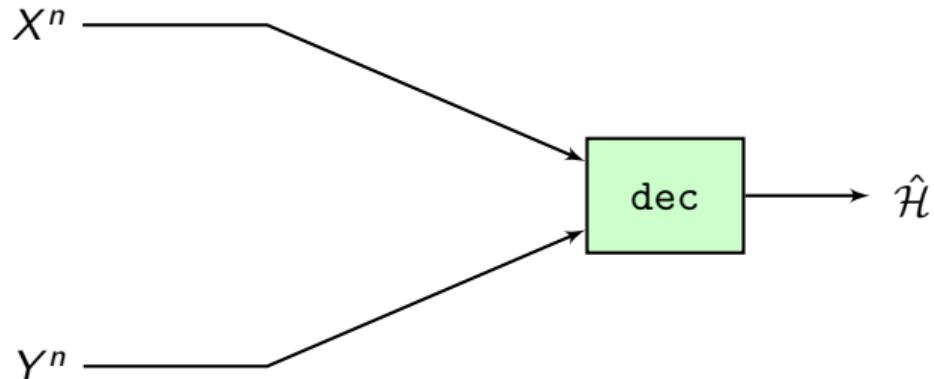
(Centralized) hypothesis testing

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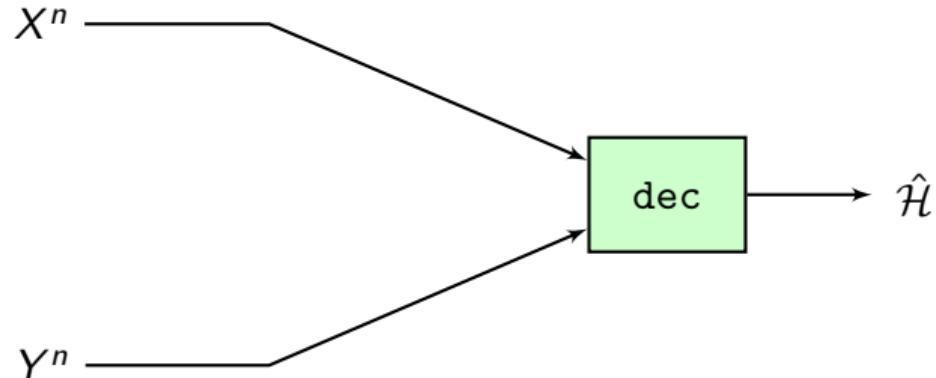
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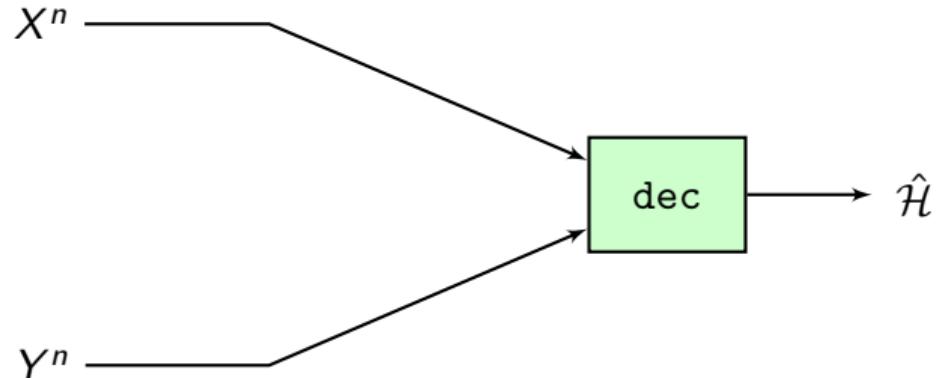
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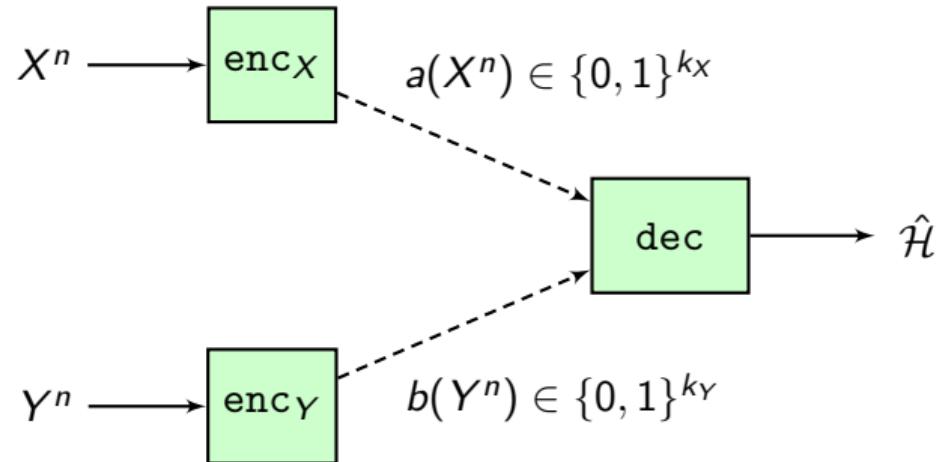
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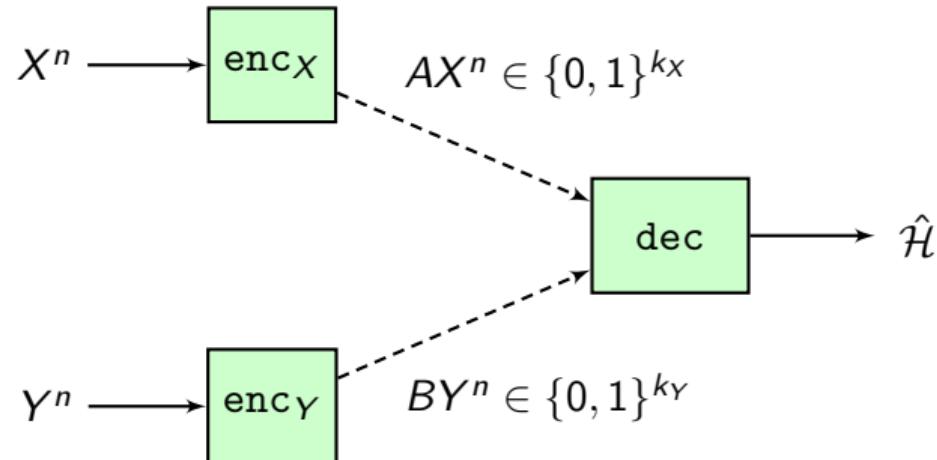
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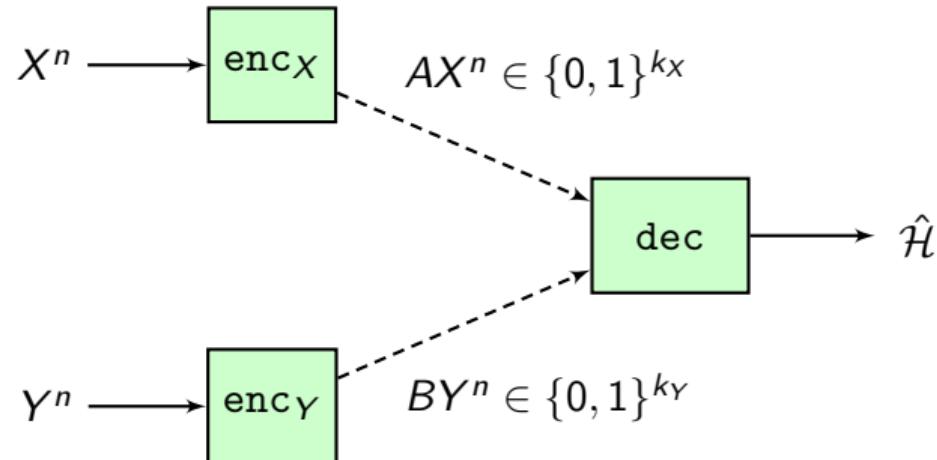
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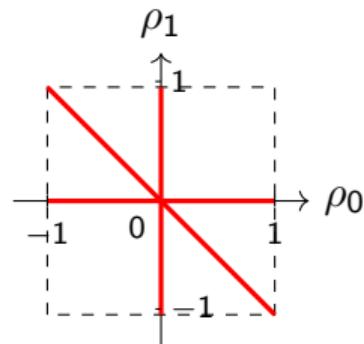
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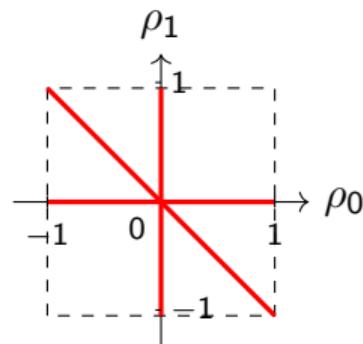
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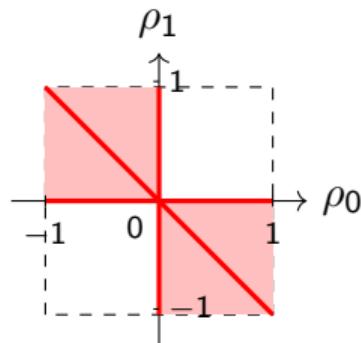
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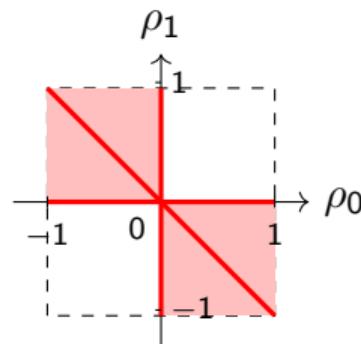
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- conjecture: $X^k \oplus Y^k \succeq AX^n \oplus AY^n$ for all ρ_0, ρ_1 of opposite signs (linear codes “bad” for testing any positive vs. negative correlation) ?

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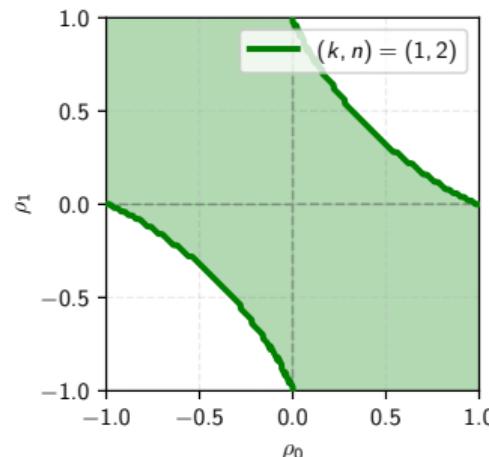
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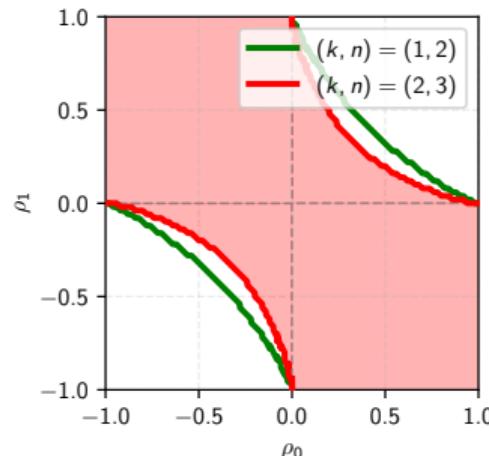
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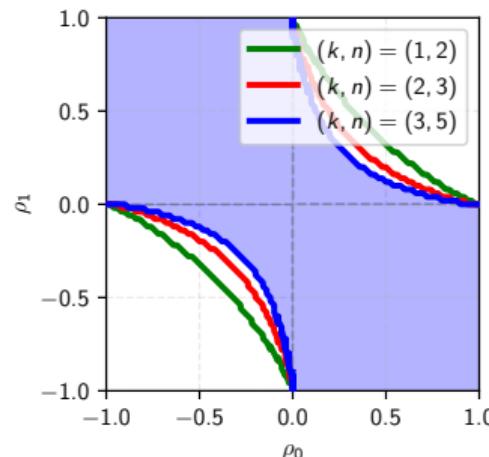
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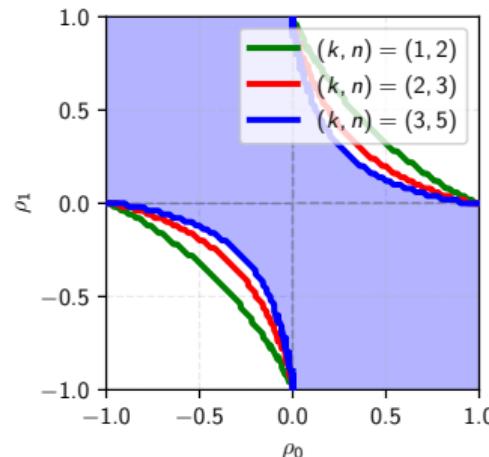
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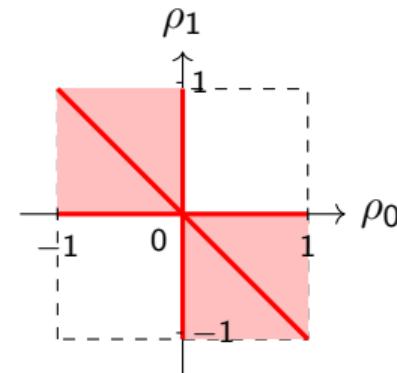
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arXiv:2601.10526

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