

Trivial is (Sometimes) Best: Distributed Hypothesis Testing via Linear Codes

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joint work with Robinson Cung and Emre Telatar

EPFL



January 30, 2026
LIDS student conference

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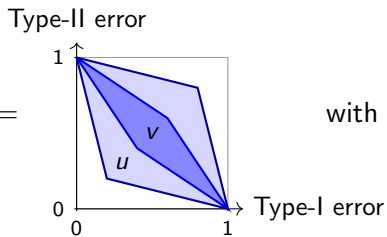
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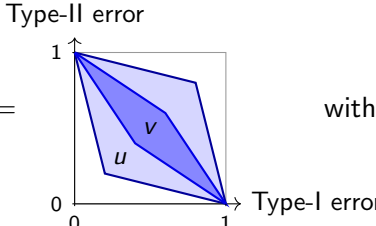
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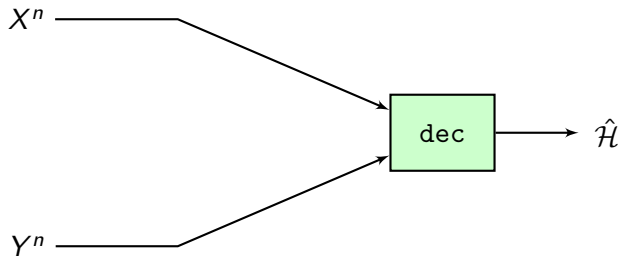
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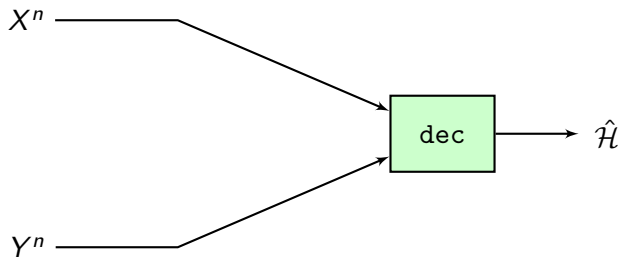
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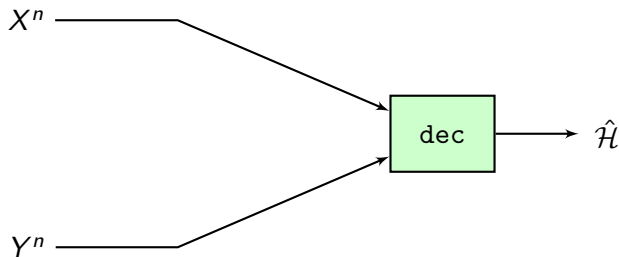
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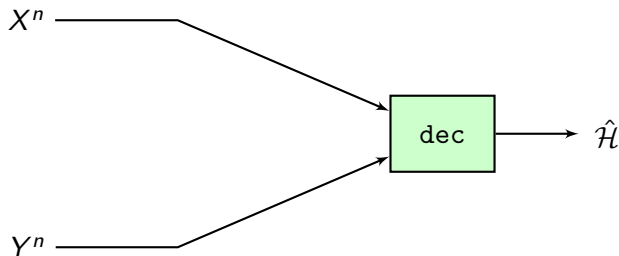
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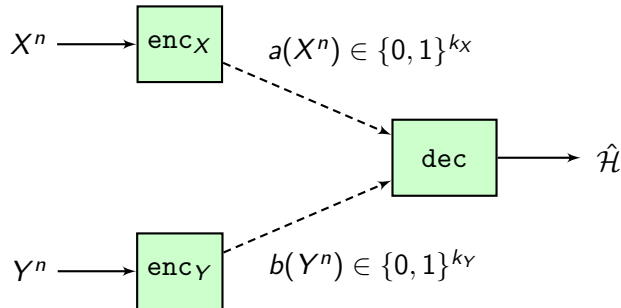
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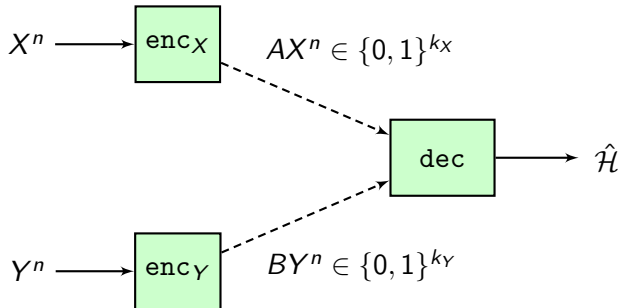
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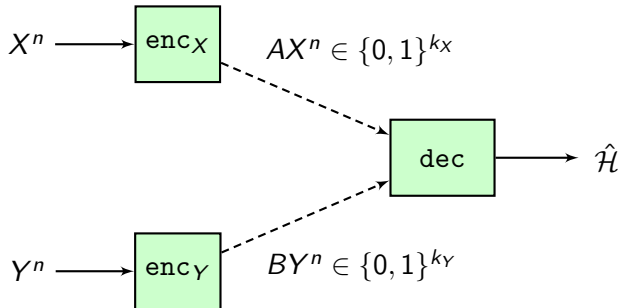
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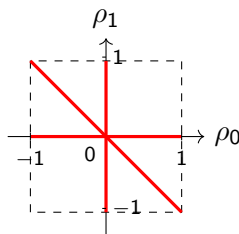
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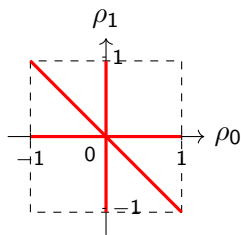
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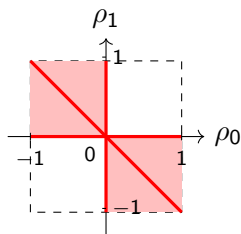
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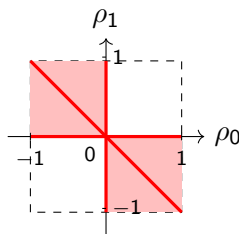
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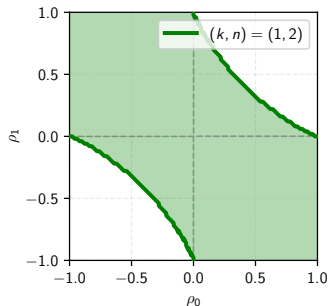
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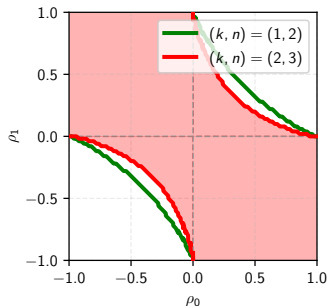
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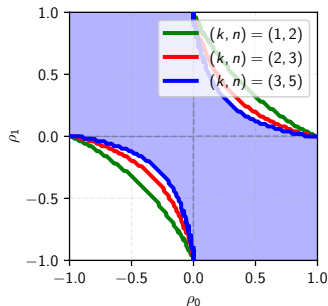
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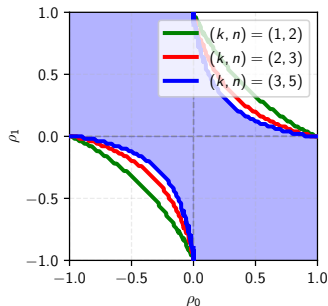
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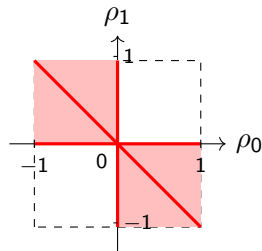
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arXiv:2601.10526

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