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IE 617: Online Learning and Bandit Algorithms Course Project Communication-Constrained Multi-Armed Bandits

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Outline



2 Theorems and Proofs

3 Simulations



Applications of Learning to Communication

- Beam alignment (Vutha Va, Takayuki Shimizu, Gaurav Bansal, et al. "Online Learning for Position-Aided Millimeter Wave Beam Training". In: *IEEE Access* (2019), Matthew B. Booth, Vinayak Suresh, Nicolò Michelusi, et al. "Multi-Armed Bandit Beam Alignment and Tracking for Mobile Millimeter Wave Communications". In: *IEEE Communications Letters* 7 (2019))
- Rate selection (Harsh Gupta, Atilla Eryilmaz, and R. Srikant. "Link Rate Selection using Constrained Thompson Sampling". In: IEEE INFOCOM 2019 - IEEE Conference on Computer Communications. 2019)
- Bit-constrained communication (Osama A. Hanna, Lin F. Yang, and Christina Fragouli. Solving Multi-Arm Bandit Using a Few Bits of Communication. 2021, Aritra Mitra, Hamed Hassani, and George J Pappas. "Linear Stochastic Bandits over a Bit-Constrained Channel". In: arXiv preprint arXiv:2203.01198 (2022))

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Problem Setup			



Source: Osama A. Hanna, Lin F. Yang, and Christina Fragouli. Solving Multi-Arm Bandit Using a Few Bits of Communication. 2021

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Problem Statement

- MAB problem, horizon n
- Leaner chooses $A_t \in A_t$ and receives r_t with mean μ_{A_t}
- Goal: maximize expected regret, $R_n = \mathbb{E}[\sum_{t=1}^n (\mu_t^* r_t)]$, where $\mu_t^* = \max_{A \in \mathcal{A}_t} \mu_A$



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Recall

Let n be the number of rounds.

- ETC and ϵ -greedy achieves $\mathcal{O}(\sqrt{n})$ with knowledge of Δ
- Thompson sampling and UCB achieves $\mathcal{O}(\sqrt{n\log n})$ without knowing Δ
- LinUCB achieves $\mathcal{O}(d\sqrt{n}\log n)$

These assumed full-precision rewards.

Goal: Develop quantization scheme to apply over *any* MAB algorithm such that the quantized regret is only a constant factor off, while maintaining a low number of bits

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Quantization

 $\mathcal{L} \text{: countable set}$

Quantizer consists of:

- $\mathcal{E}: \mathbb{R} \to \mathcal{L}$
- $\mathcal{D}: \mathcal{L} \to \mathbb{R}$

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Stochastic Quantization

Let
$$\mathcal{L} = \{\ell_i\}_{i=1}^{2^B}$$
, $x \in [\ell_1, \ell_{2^B}]$.
• $i(x) = \max\{j \mid \ell_j \le x \text{ and } j < 2^B\}$
• $\mathcal{E}_{\mathcal{L}}(x) = \begin{cases} i(x) & \text{with probability } \frac{\ell_{i(x)+1}-x}{\ell_{i(x)+1}-\ell_{i(x)}} \\ i(x)+1 & \text{with probability } \frac{x-\ell_{i(x)}}{\ell_{i(x)+1}-\ell_{i(x)}} \end{cases}$
• $D_{\mathcal{L}}(j) = \ell_j, j \in \{1, ..., 2^B\}$

Conditioned on A_t , unbiased estimate of μ_{A_t} is communicated.

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QUBAN

- Maintains Markov property, unbiasedness, bounded variance for quantized rewards
- Uses a few bits for communication

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QuBan : Main Ideas

- Center quantization scheme around value believed to be closest to picked arm's mean in majority of iterations
- Quantization error conditionally independent on past history given A_t
- Assign shorter codes to values near quantization centre and o.w. longer codes
- Use SQ to convey unbiased estimate of reward

QUBAN: Algorithm (Learner)

Algorithm 1 Learner operation with input MAB algorithm Λ

1: Initialize: $\hat{\mu}(1) = 0$

- 2: for t = 1, ..., n do
- 3: Choose an action A_t based on the bandit
- 4: algorithm Λ and ask the next agent to play it
- 5: Send M_{i} , $\hat{\mu}(t)$ to an agent
- 6: Receive the encoded reward $(b_t, I_t, \mathcal{E}_{\mathcal{L}_t}(e_t))$ (see
- 7: Algorithm 2)
- 8: Decode \hat{r}_t :
- 9: if $length(b_t) \le 4$ then
- 10: \hat{r}_t can be decoded using a lookup table
- 11: else
- 12: Decode the sign, s_t , of r_t from b_t
- 13: Set ℓ_t to be the I_t -th element in the set
- 14: $\{0, 2^0, ...\}$
- 15: Set $\mathcal{L}_t = \{\ell_t, \ell_t + 1, ..., \max\{2\ell_t, \ell_t + 1\}\}$
- 16: Let $e_t^{(q)} = D_{\mathcal{L}_t}(\mathcal{E}_{\mathcal{L}_t}(e_t))$
- 17: $\hat{r}_t = (s_t(e_t^{(q)} + \ell_t + 3.5) + 0.5 + \lfloor \hat{\mu}(t) / M_t \rfloor)M_t$
- 18: Calculate $\hat{\mu}(t+1)$ (using one of the discussed
- 19: choices)
- 20: Update the parameters required by Λ

Source: Osama A. Hanna, Lin F. Yang, and Christina Fragouli. Solving Multi-Arm Bandit Using a Few Bits of Communication. 2021 11/24

QUBAN: Algorithm (Agent)

Algorithm 2 Distributed Agent Operation

- 1: Inputs: r_t , $\hat{\mu}(t)$ and M_t
- 2: Set $L = \{\lfloor \bar{r}_t \rfloor, \lceil \bar{r} \rceil\}, \hat{\bar{r}}_t = D_L(\mathcal{E}_L(\bar{r}_t))$
- 3: Set b_t with three bits to distinguish between the 8 cases: $\hat{r}_t < -2, \hat{r}_t > 3, \hat{r}_t = i, i \in \{-1, 0, 1, 2\}.$
- 4: if $|\hat{r}_t| > |a|$ and $\hat{r}_t a > 0$, $a \in \{-2, 3\}$ then
- 5: Augment b_t with an extra one bit to indicate if $|\hat{r}_t| = |a| + 1$ or $|\hat{r}_t| > |a| + 1$.

6: **if**
$$|\hat{r}_t| > |a| + 1$$
 then

7: Let
$$L' = \{0, 2^0, ...\}$$

8: Set
$$\ell_t = \max\{j \in L | j \le |\bar{r}_t| - |a|\}$$

- 9: Encode ℓ_t by $I_t 1$ zeros followed by a one
- 10: (unary coding), where I_t is the index of ℓ_t
- 11: in the set L'.

12: Let
$$e_t = |\bar{r}_t| - |a| - \ell_t$$

13: Set
$$\mathcal{L}_t = \{\ell_t, \ell_t + 1, ..., \max\{2\ell_t, \ell_t + 1\}\}$$

14: Encode e_t using SQ to get $\mathcal{E}_{\mathcal{L}_t}(e_t)$

15: Transmit $(b_t, I_t, \mathcal{E}_{\mathcal{L}_t}(e_t))$

Source: Osama A. Hanna, Lin F. Yang, and Christina Fragouli. Solving Multi-Arm Bandit Using a Few Bits of Communication. 2021

Assumptions on MAB Instance and Algorithm

Assumption 1

All codes are prefix-free codes. Further,

- rewards possess Markov property; and
- **2** the expected regret is upper-bounded by R_n^U .

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Regret Bound

Proposition 1

Suppose Assumption 1 holds. Then, when we apply $\rm QUBAN,$ the following hold:

- Conditioned on A_t , the quantized reward \hat{r}_t is $\left(\left(1+\frac{\epsilon}{2}\right)\sigma\right)^2$ -subgaussian, conditionally independent on the history $A_1, \hat{r}_1, \ldots, A_{t-1}, \hat{r}_{t-1}$ (Markov property), and satisfies $\mathbb{E}\left[\hat{r}_t \mid A_t\right] = \mu_{A_t}, |\hat{r}_t r_t| \leq M_t$ almost surely $(t = 1, \ldots, n)$.
- **2** The expected regret R_n is bounded as $R_n \leq (1 + \frac{\epsilon}{2}) R_n^U$, where ϵ is a parameter to control the regret vs number of bits trade-off.

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Number of Bits

Theorem 1

Suppose Assumption 1 holds. Let $\epsilon = 1$. There is a universal constant C such that, for QUBAN with:

- $\hat{\mu}(t) = \hat{\mu}_{A_t}(t-1)$ (avg-arm-pt), the average number of bits communicated satisfies that $\mathbb{E}[\bar{B}(n)] \leq 3.4 + \frac{c}{n} \sum_{i=1}^{k} \log(1 + |\mu_i|/\sigma) + C/\sqrt{n}.$

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Lower Bound

Theorem 2

For any memoryless algorithm that only uses quantized rewards, prefix-free encoding and satisfies that for any MAB instance with subgaussian rewards:

R_n is sublinear in n,

2 Conditioned on r_t , $\hat{r}_t - r_t$ is $\left(\frac{\sigma}{2}\right)^2$ -subgaussian (t = 1, ..., n),

there exist σ^2 -subgaussian reward distributions for which:

- $(\forall b \in \mathbb{N})(\exists t, \delta > 0)$ such that $\mathbb{P}[B_t > b] > \delta$.
- $(\forall t > 0) (\exists n > t) \text{ such that } \mathbb{E}[\overline{B}(n)] \ge 2.2 \text{ bits.}$

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Upper Bound

$$\begin{split} \mathcal{B}_{t} &\leq 3+\mathbf{1} \left[\frac{r_{t}}{M_{t}} - \left\lfloor \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor > 3 \right] + \mathbf{1} \left[\left\lfloor \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor - \frac{r_{t}}{M_{t}} > 2 \right] \\ &+ 2 \left(\mathbf{1} \left[\frac{r_{t}}{M_{t}} - \left\lfloor \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor > 4 \right] + \log \left(\frac{r_{t}}{M_{t}} - \left\lfloor \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor - 3 \right) \right] \right) \\ &+ 2 \left(\mathbf{1} \left[\left\lfloor \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor - \frac{r_{t}}{M_{t}} > 3 \right] \left[\log \left(\left\lfloor \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor - \frac{r_{t}}{M_{t}} - 2 \right) \right] \right) \mathcal{B}_{t} \leq 3 + \\ &+ 2 \left(\mathbf{1} \left[\left\lfloor \frac{r_{t}}{M_{t}} - \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor > 3 \right] \log \left(\left\lfloor \frac{r_{t}}{M_{t}} - \frac{\hat{\mu}(t)}{M_{t}} \right\rfloor - 2 \right) \right) \mathcal{B}_{t} \leq 3 + \\ &+ 2 \left(\mathbf{1} \left[\left\lfloor \frac{r_{t} - \mu_{A_{t}}}{\sigma} \right\rfloor > 3(1 - \delta) \right] + \mathbf{1} \left[\left\lfloor \frac{\mu_{A_{t}} - \hat{\mu}(t)}{\sigma} \right\rfloor > 3\delta \right] \right) \\ &+ 2 \left(\mathbf{1} \left[\left\lfloor \frac{r_{t} - \mu_{A_{t}}}{\sigma} \right\rfloor > 3 \right] \right) \log \left(\left\lfloor \frac{r_{t} - \hat{\mu}(t)}{\sigma} \right\rfloor - 2 \right) \text{ for each } \delta > 0 \end{split}$$

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Upper Bound

$$\mathbb{E}\left[B_{t}\right] \leq 3 + \mathbb{P}\left[\left|\frac{r_{t} - \mu_{A_{t}}}{\sigma}\right| > 2(1 - \delta)\right] + \mathbb{P}\left[\left|\frac{\mu_{A_{t}} - \hat{\mu}(t)}{\sigma}\right| > 2\delta\right] \\ + 2\left(\mathbb{P}\left[\left|\frac{r_{t} - \mu_{A_{t}}}{\sigma}\right| > 3(1 - \delta)\right] + \mathbb{P}\left[\left|\frac{\mu_{A_{t}} - \hat{\mu}(t)}{\sigma}\right| > 3\delta\right]\right) \\ + 2\mathbb{E}\left[\left(1\left[\left|\frac{r_{t} - \mu_{A_{t}}}{\sigma}\right| > 3\right]\right)\log\left(\left|\frac{r_{t} - \hat{\mu}(t)}{\sigma}\right| - 2\right)\right] \\ \leq 3.4 + C\mathbb{E}\left[\left|\frac{\mu_{A_{t}} - \hat{\mu}(t)}{\sigma}\right|\right] \leq \cdots \quad \Box$$

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QUBAN



Pseudoregret vs. T for QUBAN

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Modified setup

- Agent full precision, learner bit-constrained? Trivial.
- Both bit-constrained?

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${\small {\sf Modified}} \ {\rm QUBAN}$

- Learner too is communication-constrained
- Learner sends $\hat{\mu}(t)$ using 10-bit SQ

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Modified QUBAN



Pseudoregret vs. T for modified QUBAN

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Conclusion

- Presented the upper bound proof, and
- Numerical analysis for the setup where both the learner and agents are bit-constrained.

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Thank you